

Domination Numbers Of Grids $P_{14} \times P_n$
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□ **ABSTRACT** □

This paper concerns the domination numbers $\gamma(P_k \times P_n)$ for the $P_k \times P_n$ grid graphs for $k = 14$ and for all $n \geq 1$.

These numbers were previously established for $1 \leq n \leq 12$ [2], [5].

Keywords: dominating Set, domination number, transformation of a dominating set, Cartesian product of two paths.

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مراتب السيطرة للشبكات $P_{14} \times P_n$ 05C69 تصنيف موضوع الرياضيات

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□ ملخص □

هذه الورقة تهتم بمراتب السيطرة $\gamma(p_k \times p_n)$ لأجل البيانات الشبكية $p_k \times p_n$.
من أجل $k = 14$ ومن أجل كل $n \geq 1$.
هذه المراتب درست من قبل لأجل $[5]; [2]; 1 \leq n \leq 12$.

الكلمات المفتاحية: مجموعة سيطرة، مرتبة سيطرة، نقل مجموعة سيطرة، الجداء الديكارتي لمسارين.

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Introduction:

A dominating set in a graph is a set of vertices having the property that every vertex not in the set is adjacent to a vertex in the set. The domination number $\gamma(G)$ of a graph G is the cardinality of a smallest dominating set in G .

The problem of finding the domination number of an arbitrary grid graph (=subgraph of $P_k \times P_n$) is NP-complet [3], [4].

In this paper, we introduce the concept of transforming the domination

From a vertex in a dominating set D of graph $G = (V, E)$ to a vertex in

$V - D$, where G is a simple connected graph. We give an algorithm using this transformation to obtain a dominating set of a graph G .

A graph $G = (V, E)$ is a mathematical structure which consists of two sets V and E , where V is finite and nonempty, and every element of E is an unordered pair $\{u, v\}$ of distinct elements of V ; we simply write uv instead of $\{u, v\}$.

The elements of V are called vertices, while the elements of E are called edges [1].

Two vertices u and v of a graph G are said to be adjacent if $uv \in E$.

The neighborhood of v is the set of all vertices of G which are adjacent to v ;

the neighborhood of v is denoted by $N(v)$. The closed neighborhood of v is $\bar{N}(v)$, $\bar{N}(v) = N(v) \cup \{v\}$.

The degree $d(v)$ of a vertex v is the cardinality $|N(v)|$, $d(v) = |N(v)|$.

Definition:

Let D be a dominating set of a graph $G = (V, E)$.

1. We define the function C_D , which we call the weight function, as follows: $C_D : V \rightarrow \mathbb{N}$, where \mathbb{N} is the set of natural numbers, $C_D(v) = |\bar{N}(v)|$,

where $\bar{N}(v) = \{w \in D : vw \in E \text{ or } w = v\}$

i.e. the weight of v is the number of vertices in D which dominate v .

2. We say that $v \in D$ has a moving domination if there exists a vertex $w \in \bar{N}(v) - D$ such that $wu \in E$ for every vertex $u, u \in \{x \in N(v) : C_D(x) = 1\}$.

3. We say that a vertex $v \in D$ is a redundant vertex of D if $C_D\{w\} \geq 2$ for every vertex $w \in \bar{N}(v)$.

4. If $v \in D$ has a moving domination, we say that v is inefficient if transforming the domination from v to any vertex in $\bar{N}(v)$ would not produce any redundant vertex.

Grid graph $P_k \times P_n$:

For two vertices v_0 and v_n of a graph G , a $v_0 - v_n$ walk is an alternating sequence of vertices and edges $v_0, e_1, v_1, \dots, e_n, v_n$ such that consecutive vertices and edges are incident.

A path is a walk in which no vertex is repeated. A path with n vertices is denoted by P_n , it has $n - 1$ edges; the length of P_n is $n - 1$; the cartesian product $P_k \times P_n$ of two paths is the grid graph with vertex set

$$V = \{(i, j) : 1 \leq i \leq k, 1 \leq j \leq n\}$$

where $(u_1, u_2)(v_1, v_2)$ is an edge of $P_k \times P_n$ if $|u_1 - v_1| + |u_2 - v_2| = 1$ [4].

If D is a dominating set of $P_k \times P_n$ which has no redundant vertex, then a vertex $v \in D$ has a moving domination if and only if one of the following two cases occurs:

Case (1): for every vertex $w \in \bar{N}(v)$, we have $C_D(w) \geq 2$.

In this case, the domination of v can be transformed to any vertex in $N(v) - D$.

Case(2): there exists exactly one vertex $u \in N(v)$ such that $C_D(u) = 1$. In this case, the domination of v can be transformed only to u .

Algorithm for finding a dominating set of a graph $P_k \times P_n$ using a transformation of domination of vertices:

1. Let $P_k \times P_n = (V, E)$ be a graph of order greater than 1; $|V| = m$.
 2. Let $D = V$ be a dominating set of $P_k \times P_n$.
then, for any vertex $v \in D$ we have $C_D(v) = d(v) + 1 \geq 2$.
 3. Pick a vertex v_1 of D , and delete from D all vertices $w, w \in N(v_1)$.
then, for $1 < n < m/2$, pick a vertex $v_n \in D - \bigcup_{i=1}^{n-1} \bar{N}(v_i)$ and delete from D all vertices $w, w \in N(v_n) - \bigcup_{i=1}^{n-1} \bar{N}(v_i)$. If D contains a redundant vertex, then delete it.
Repeat this process until D has no redundant vertex.
 4. Transform domination from vertices of D which have moving domination to vertices in $V - D$ to obtain redundant vertices and go to step 4.
- If no redundant vertex can be obtained by a transformation of domination of vertices of D , then stop, and the obtained dominating set D satisfies:

for every $v \in D, \exists w \in \bar{N}(v)$ such that $C_D(w) = 1$.

Example:

1. Let (k, n) be the vertex in the k -th row and in the n -th column of the graph $G = P_{14} \times P_{12}; |V| = 168$.
2. Let $D = V$, dominating set of G .
3. Pick a vertex $v_1 = (1, 2) \in D$, and delete from D all vertices $w, w \in N(v_1)$, then, for $1 < n < 168/2$, pick a vertex $v_n, v_n \in D - \bigcup_{i=1}^{n-1} \bar{N}(v_i)$, and delete from D all vertices $w, w \in N(v_n) - \bigcup_{i=1}^{n-1} \bar{N}(v_i)$. We obtain the dominating set D (black circles) in figure 1.

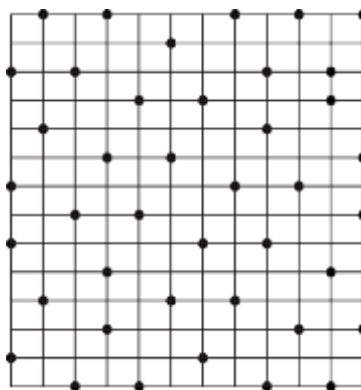


Figure1.

4. Since for every vertex $v \in D, \exists w \in \bar{N}(v)$ such that $C_D(w) = 1$, D has

no red un dantvertices.

5. Transform the domination from the vertex (1, 12) to the vertex (2, 12) and delete, from D , the resulting redundant vertex (3, 11).

The set D indicated in figure 2 (black circles) is a dominating set of $G = P_{14} \times P_{12}$

Note that D is a minimum dominating set (see [2]). $\gamma(P_{14} \times P_{12}) = 40$.

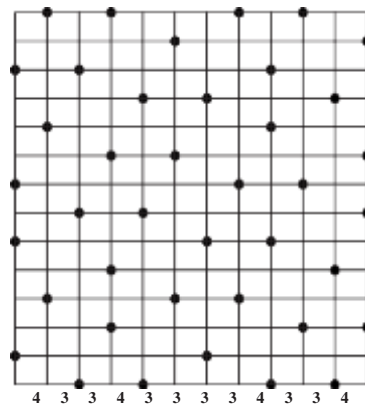


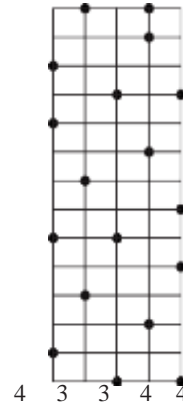
Figure2.

$$n = 12, \gamma(P_{14} \times P_{12}) = 4(n - 8) + 24 = 4n - 8$$

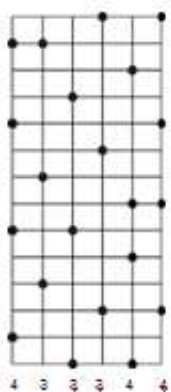
By the same method, we give $\gamma(P_{14} \times P_n); n \geq 1$.



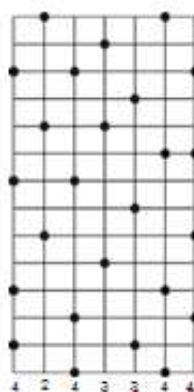
$$n = 4, \gamma(P_{14} \times P_4) = 4(n - 2) + 6 = 4n - 2$$



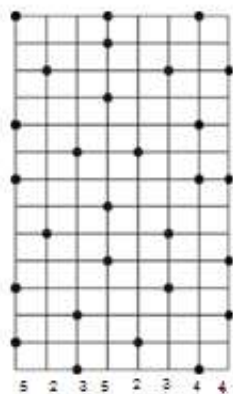
$$n = 5, \gamma(P_{14} \times P_5) = 4(n - 2) + 6 = 4n - 2$$



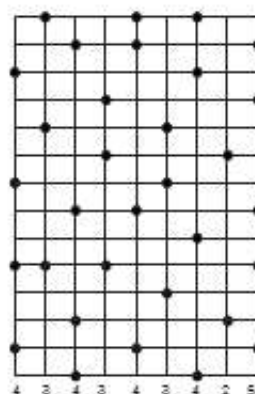
$$n = 6, \gamma(P_{14} \times P_6) = 4(n-3) + 9 \\ = 4n - 3$$



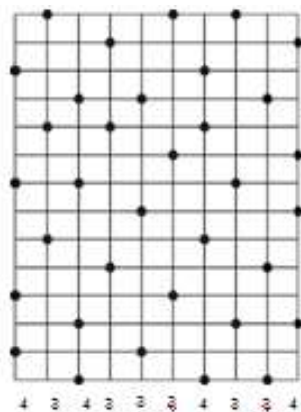
$$n = 7, \gamma(P_{14} \times P_7) = 4(n-3) + 8 \\ = 4n - 4$$



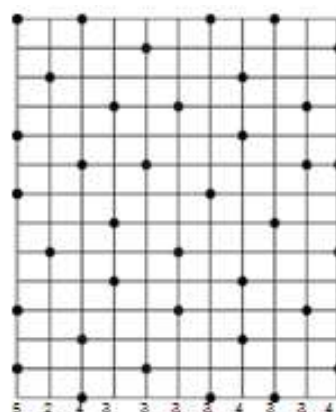
$$n = 8, \gamma(P_{14} \times P_8) = 4(n-6) + 20 \\ = 4n - 4$$



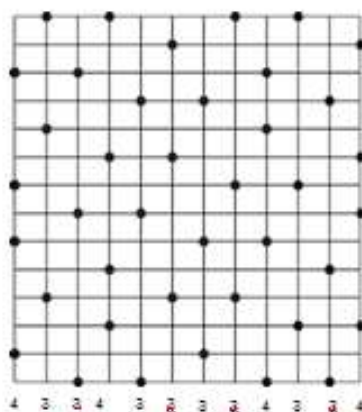
$$n = 9, \gamma(P_{14} \times P_9) = 4(n-5) + 16 \\ = 4n - 5$$



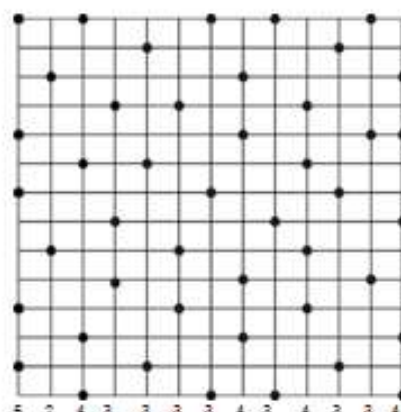
$$n = 10, \gamma(P_{14} \times P_{10}) = 4(n-6) + 18 = 4n - 6$$



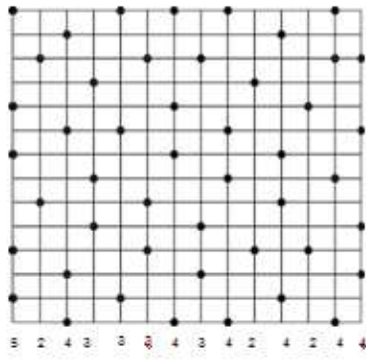
$$n = 11, \gamma(P_{14} \times P_{11}) = 4(n-8) + 25 = 4n - 7$$



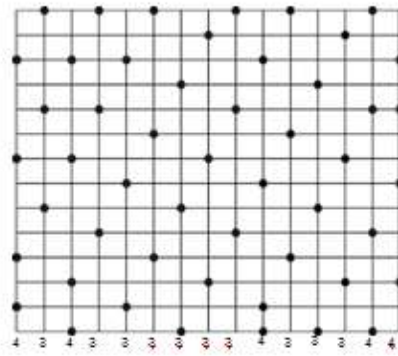
$$n = 12, \gamma(P_{14} \times P_{12}) = 4(n-8) + 24 = 4n - 8$$



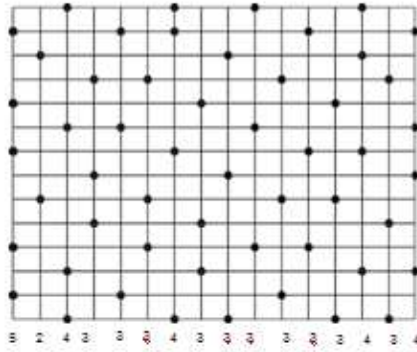
$$n = 13, \gamma(P_{14} \times P_{13}) = 4(n-9) + 28 = 4n - 8$$



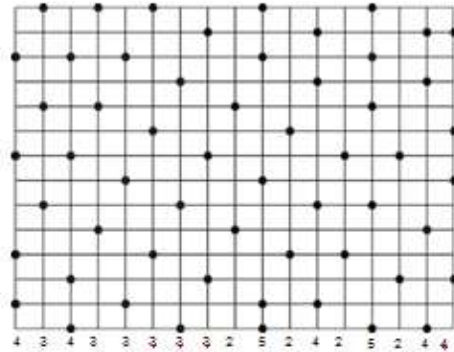
$$n = 14, \gamma(P_{14} \times P_{14}) = 4(n - 8) + 23 = 4n - 9$$



$$n = 15, \gamma(P_{14} \times P_{15}) = 4(n - 10) + 30 = 4n - 10$$



$$n = 16, \gamma(P_{14} \times P_{16}) = 4(n - 12) + 17 = 4n - 11$$



$$n = 17, \gamma(P_{14} \times P_{17}) = 4(n - 12) + 36 = 4n - 12$$

Hence :

$$\gamma(P_{14} \times P_n) = \begin{cases} 5 ; n = 1 \\ 8 ; n = 2 \\ 11 ; n = 3 \\ 14 ; n = 4 \\ 18 ; n = 5 \\ 21 ; n = 6 \\ 4n - (4 + 8t + 18k) ; n = 7 + 10t + 22k \\ 4n - (m + 4t + 18k) ; n = 4 + m + 5t + 22k, 4 \leq m \leq 7. \\ 4n - (8 + 9t + 18k) ; n = 12 + 11t + 22k \\ 4n - (m + 5t + 18k) ; n = 6 + m + 6t + 22k, 12 \leq m \leq 16. \end{cases}$$

where t, k and m are positive integers, $0 \leq t \leq 1, k \geq 0$.

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